# SMS++: a Structured Modelling System for Optimization

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PGMO Days 2019 December 4, 2019, EDF'Lab Palaiseau, France







#### **Outline**

• The plan4res project

- SMS++
- Stochastic SMS++
- Current and future work

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### The plan4res H2020 project

- "An end-to-end planning and operation tool, composed of a set of optimization models based on an integrated modelling of the European Energy System"
- ullet An accurate depiction of long-term effects of strategic choices on the pan-European Energy System  $\equiv$

modelling the next 30 years with 1h timescale and huge amounts of uncertainty over everything

An unfeasibly large optimization problem with lots of structure

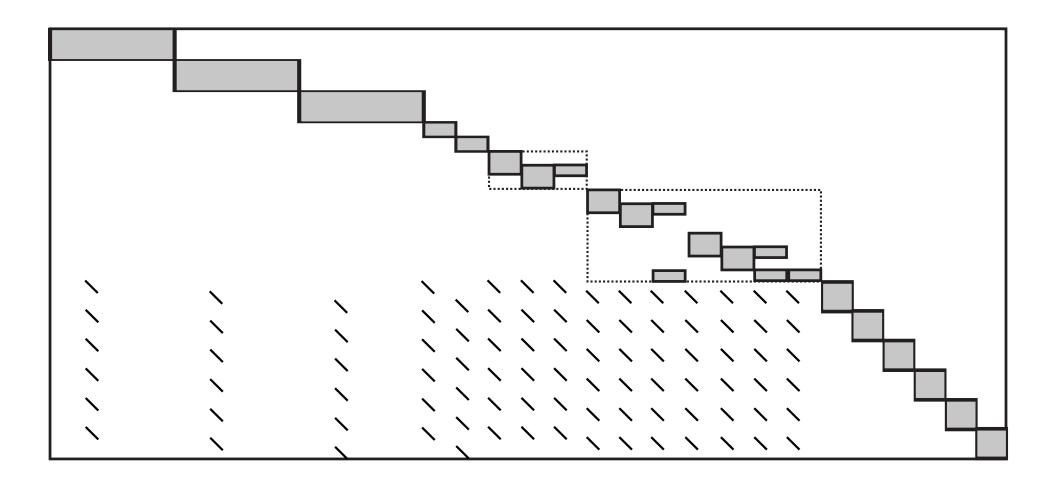
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## **Short-term problem: Unit Commitment (UC)**

- Find a (near-)optimal schedule of a large number of units satisfying the demand at each node of the network, while respecting a set of technical constraints, at each time instant of the horizon (e.g., 1 hour)
- Three natural sources of structure: unit, time, and network
- Relaxing demand constraints decomposes by unit and network: one problem per unit across all horizon and a network problem per time instant

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## Short-term problem: Unit Commitment (UC)



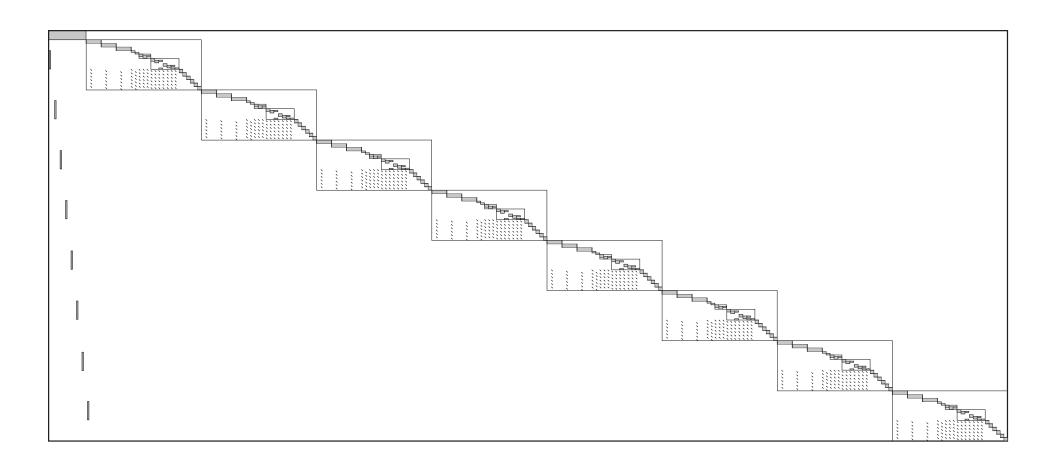
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## Mid-term problem: Seasonal storage valuation

- UC is a (deterministic) short-term problem and lacks long-term strategies
- The mid-term (e.g., 1 year) problem provides the UC with approximations of the cost-2-go function
- UC then arises at each stage (e.g., 1 week) of the mid-term problem
- Uncertainties: inflows, demand, outages, intermittent generation
- A multi-stage stochastic optimization problem

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## Mid-term problem: Seasonal storage valuation



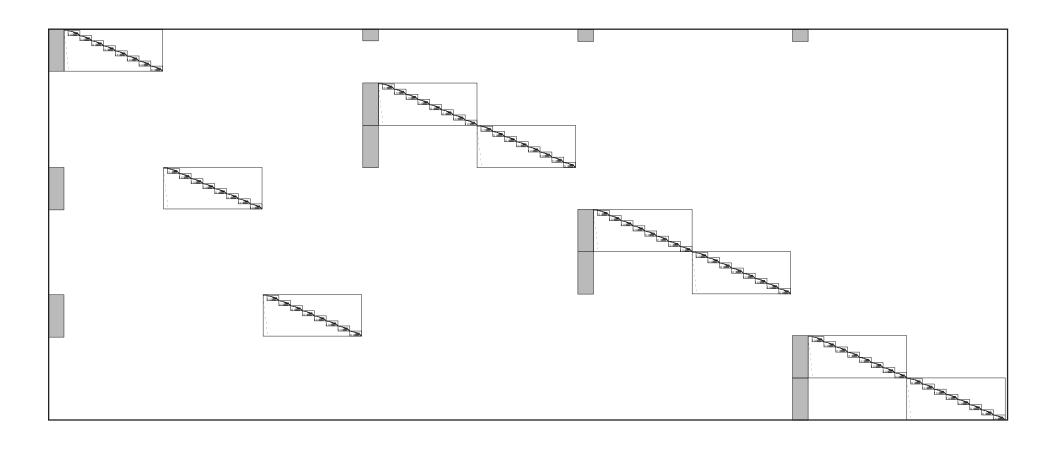
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### Long-term problem: Investment layer

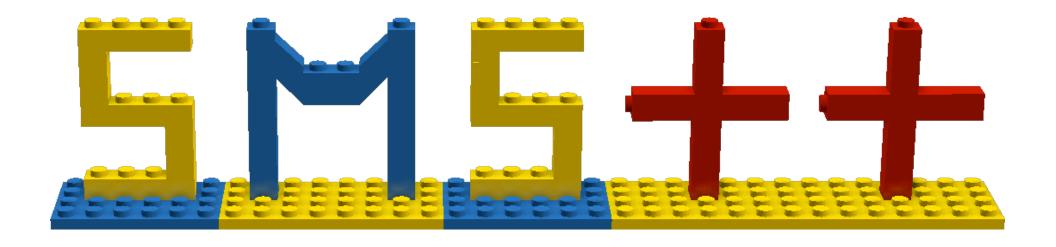
- Long-term planning needed: the energy system changes frequently,
   but modifications are slow and costly
- Uncertainties in demand and production:
  - shifts in consumption patterns (EV, cryptocurrencies, . . . )
  - regulatory factors (EU energy market, . . . ),
  - political factors ( $CO_2$  emission treaties, nuclear power, . . . )
  - **–** . . .
- Design the optimal generation mix with the optimal transmission and distribution grid capacities
- 30-year horizon with 5-year steps (multi-level recourse), many scenarios

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## Long-term problem



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#### SMS++

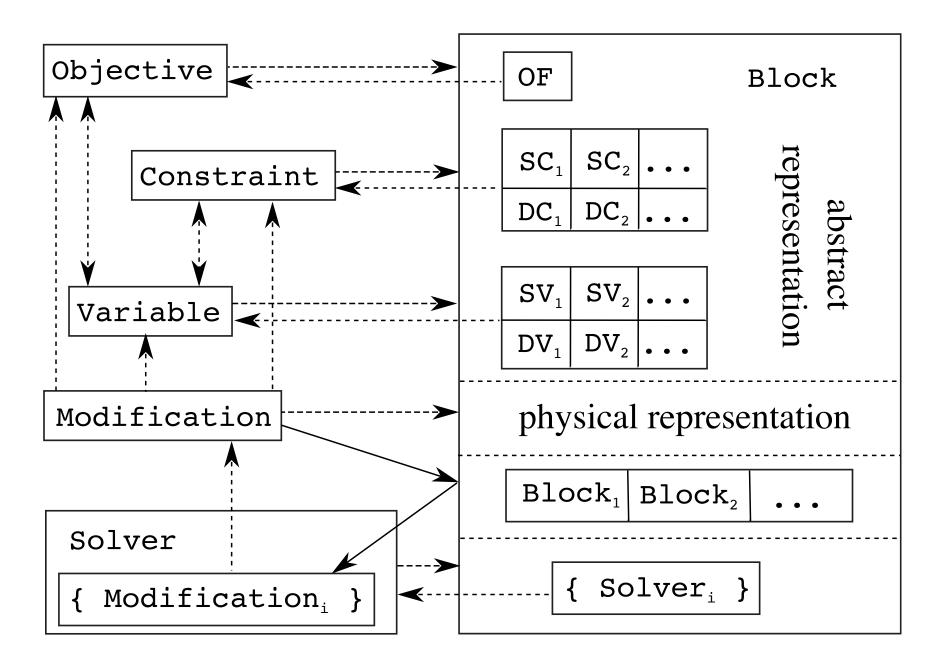
A set of C++ classes implementing a modelling system that:

- explicitly supports the notion of block ≡ nested structure
- separately provides "semantic" information from "syntactic" details (objective and list of constraints/variables 

   one specific formulation among many)
- allows exploiting specialised solvers on blocks with specific structure
- manages dynamic changes in the model beyond "just" generation of constraints/variables
- manages reformulation/restriction/relaxation

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#### SMS++



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- Block is an abstract class representing the general concept of "a part of a mathematical model with a well-understood identity".
- Each :Block is a model with specific structure
   (e.g., MCFBlock:Block = a Min-Cost Flow problem).
- Physical representation of a Block: whatever data structure is required to describe the instance (e.g., for a MCFBlock, a graph, source and sink nodes, cost and capacity of each arc, . . . )
- Abstract representation of a Block: an Objective and an "unstructured" list of Constraints and Variables.

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- Solver for a specific :Block can use the physical representation
   no need for explicit Constraint or Objective
   abstract representation of Block only constructed on demand
- A general-purpose Solver uses the abstract representation
- Dynamic Variable/Constraint can be generated on demand (user cuts/lazy constraints/column generation)
- ullet Objective of sub-Blocks summed to that of father Block if it has same sense, otherwise  $\min/\max$

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- Solver = interface between a Block and algorithms solving it
- Each Solver attached to a single Block, from which it picks all the data, but any # of Solver can be attached to the same Block
- Individual Solver can be attached to sub-Block of a Block
- Tries to cater for all the important needs:
  - optimal and sub-optimal solutions, provably unbounded/infeasible
  - time/resource limits for solutions, but restarts (reoptimization)
  - any # of multiple solutions produced on demand
  - lazily reacts to changes in the data of the Block via Modification

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## SMS++ Block and Modification

- Most Block components can change, but not all:
  - set of sub-Block
  - # and shape of groups of Variable/Constraint
- Any change is communicated to each interested Solver (attached to the Block or any of its ancestor) via a Modification object
- However, two different kinds of Modification (what changes):
  - physical Modification, only specialized Solver concerned

abstract Modification, only Solver using it concerned

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### **Example: Capacitated Facility Location**

Given a set L of locations and a set D of customers, the problem consists in selecting a subset of the locations in which facilities will be placed in order to serve the given set of customers.

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### **Example: Capacitated Facility Location**

$$\begin{aligned} & \min & & \sum_{i \in L} f_i y_i + \sum_{i \in L} \sum_{j \in D} c_{ij} x_{ij} \\ & \text{s.t.} & & \sum_{i \in L} x_{ij} = 1, \forall j \in D \\ & & \sum_{i \in L} d_j x_{ij} \leq u_i y_i, \forall i \in L \\ & & & x \geq 0 \\ & & & y \in \{0,1\}^{|L|} \end{aligned}$$

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## **Example: Capacitated Facility Location Abstract representation**

min 
$$f(x,y)$$
  
s.t.  $h_i(x) = 0, \forall i \in I$   
 $g_j(x) \leq 0, \forall j \in J$   
 $x \in X$   
 $y \in Y$ 

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## **Example: Capacitated Facility Location Physical representation**

- ullet L and D
- $f_i, \ \forall i \in L$
- $c_{ij}, \ \forall \ i \in L, \ \forall j \in D$
- $d_j$ ,  $\forall j \in D$

## **Example: Capacitated Facility Location**

$$\min \sum_{i \in L} f_i y_i$$
 s.t. 
$$\sum_{j \in D} d_j x_{ij} \le u_i y_i, \forall i \in L$$
 
$$y \in \{0,1\}^{|L|}$$

j-th sub-Block:

$$\min \sum_{i \in L} c_{ij} x_{ij}$$
  
s.t. 
$$\sum_{i \in L} x_{ij} = 1$$
  
$$x_j \ge 0$$

SMS++ is (almost) ready for deterministic optimization.

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SMS++ is (almost) ready for deterministic optimization.

What about stochastic optimization?

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We must represent uncertainty in SMS++.

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We must represent uncertainty in SMS++.

UCBlock ThermalUnitBlock

UnitBlock EMobilityUnitBlock

HeatBlock PowerToGasUnitBlock

NetworkBlock BatteryUnitBlock

HydroUnitBlock IntermittentUnitBlock

DCNetworkBlock CentralizedDemandResponseUnitBlock

BusNetworkBlock ...

Ideally, without changing the implementation of the Blocks.

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$$\min_{\mathbf{s.t.}} f_{1}(x_{1}) + \mathbb{E} \left[ \min_{\mathbf{s.t.}} f_{2}(x_{2};\xi_{2}) + \mathbb{E}_{|\xi_{[2]}|} + \mathbb{E}_{|\xi_{[2]}|} \right] + \mathbb{E}_{|\xi_{[2]}|} \left[ \cdots + \mathbb{E}_{|\xi_{[T-1]}|} \left[ \min_{\mathbf{s.t.}} f_{T}(x_{T};\xi_{T}) \right] \right] \right]$$

 $\{\xi_t\}_{t\in\{2,...,T\}}$  is a stochastic process

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$$\min_{\mathbf{s.t.}} f_{1}(x_{1}) + \mathbb{E} \left[ \min_{\mathbf{s.t.}} f_{2}(x_{2};\xi_{2}) + \mathbb{E}_{|\xi_{[2]}|} + \mathbb{E}_{|\xi_{[2]}|} \right] + \mathbb{E}_{|\xi_{[2]}|} \left[ \cdots + \mathbb{E}_{|\xi_{[T-1]}|} \left[ \min_{\mathbf{s.t.}} f_{T}(x_{T};\xi_{T}) \right] \right] \right]$$

 $\{\xi_t\}_{t\in\{2,...,T\}}$  is a stochastic process

Stochastic dual dynamic programming (SDDP)

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$$V_{t}(x_{t-1}, \xi_{t}) = \min_{x_{t}} f_{t}(x_{t}; \xi_{t}) + \mathcal{V}_{t+1}(x_{t}, \xi_{t})$$
s.t.  $x_{t} \in X_{t}(x_{t-1}, \xi_{t})$ 

$$\mathcal{V}_{t+1}(x_t, \xi_t) = \mathbb{E}\left[V_{t+1}(x_t, \xi_{t+1}) \mid \xi_{[t]}\right]$$

$$V_t(x_{t-1}, \xi_t) = \min_{x_t} f_t(x_t; \xi_t) + \mathcal{V}_{t+1}(x_t, \xi_t)$$
  
s.t.  $x_t \in X_t(x_{t-1}, \xi_t)$ 

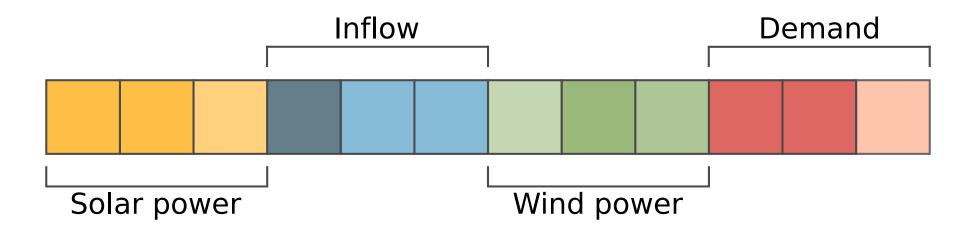
$$\mathcal{V}_{t+1}(x_t, \xi_t) = \mathbb{E}\left[V_{t+1}(x_t, \xi_{t+1}) \mid \xi_{[t]}\right]$$

$$\min_{x_t} f_t(x_t; \tilde{\xi}_t) + \mathcal{P}_{t+1}(x_t)$$
s.t.  $x_t \in X_t(\tilde{x}_{t-1}, \tilde{\xi}_t)$ 

We must be able to:

- Simulate the random variables.
- Update the data of the Blocks for a given realization of the random variables.

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### Setting the data

We can set the data by using either the abstract or the physical representation of the Block.

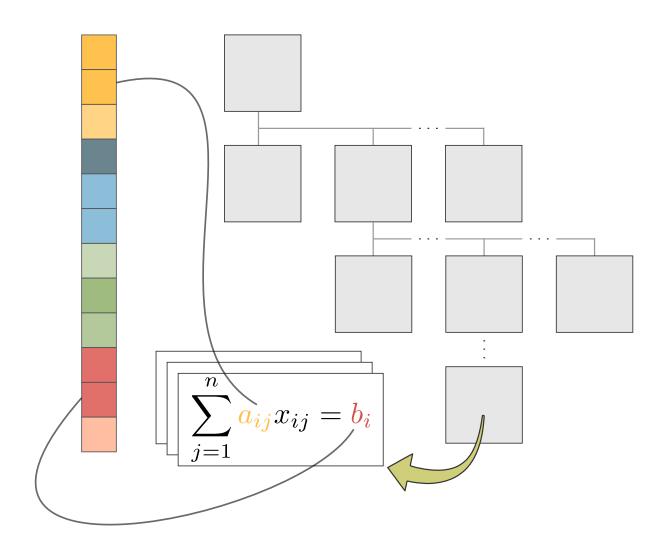
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#### **Setting the data Abstract representation**

- A random variable may appear in the Objective and/or in the Constraints.
- Let's assume that a random variable appears in the right-(or left-)hand side of a RowConstraint or as the coefficient of some Variable.

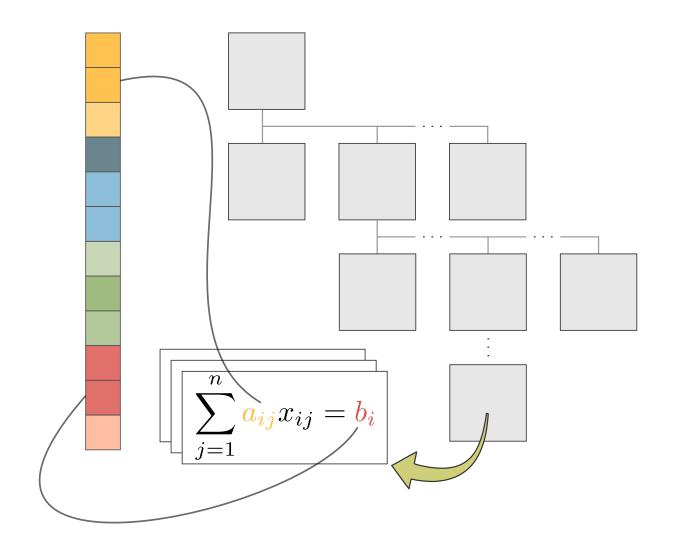
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## **Setting the data Abstract representation**



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## **Setting the data Abstract representation**



It is doable, but can be complicated for the user.

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## **Setting the data** Physical representation

 We set the data by using the available methods in the Blocks.

```
void HydroUnitBlock::set_inflow(std::vector<double> inflow);
```

- The methods are first registered in the methods factory.
- Pointer to a method can be retrieved by its name.

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## **Setting the data** Physical representation

- What if the available methods are not enough?
- One can write their own method and register it in the methods factory.

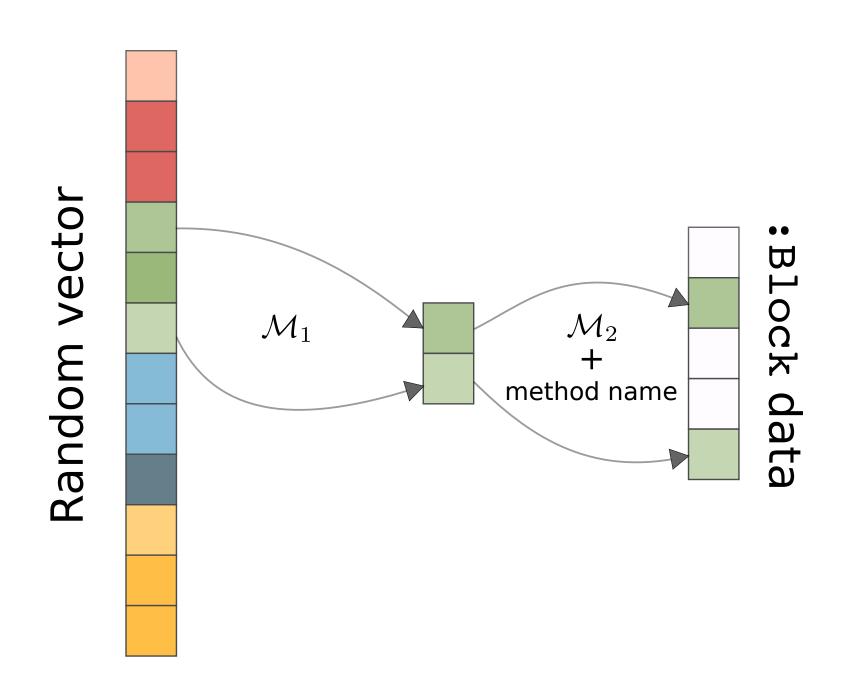
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### **Data Mapping**

- Data mapping identifies the random variables in the Blocks.
- At the same time, it provides means to set the values of those variables.
- It associates methods in the methods factory with instances of :Block.

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## **Data Mapping**



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#### StochasticBlock

• It has a (single) nested Block (which is becoming stochastic).

It has a data mapping.

• It has a probability distribution or a "partial stochastic process".

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#### **Current and future work**

- BendersBlock turns a Block into its Benders' reformulation
- LagrangianDualBlock turns any Block into its
   Lagrangian Dual w.r.t. constraints linking its sub-Block
- Asynchronous execution of the computationally heavy parts.

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### **Acknowledgements**

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This project has received funding from the European Union's Horizon 2020 research and innovation programme under grant agreement No 773897



